## Homework #4 Solutions

## Problem 1

a. The lower bound is 1. If n is even, let X be ((c, c), ..., (b, b)...), where (c, c) is played for n/2 periods and (b, b) is played for n/2 periods. For n odd, X = ((c, c), ..., (b, b), ..., (a, a)), where (c, c) is played for (n - 1)/2 periods, (b, b) is played for (n - 1)/2 periods. The nash equilibrium strategy is to play X as long as X has been played in every previous period, and otherwise play (c, c) for the rest of the game. This is a nash equilibrium because there is no possible deviation for either player. If a player deviates, he will get payoff at most 1 in every future period, so the best he can get by deviating is payoff n, which he is indifferent to.

b. For *n* even,  $X_n = ((c, c), ..., (b, b), ..., (a, a))$ , where (c, c) is played for n/2 periods, (b, b) is played for n/2-1 periods. If *n* is odd, let *X* be ((c, c), ..., (b, b)...), where (c, c) is played for (n - 1)/2 periods and (b, b) is played for (n + 1)/2 periods. For n = 1, the strategy is to play (b, b), for payoff 2. We prove by induction.

Suppose that for n < T, there is a subgame perfect equilibrium with payoff n + 1. At n = T, the subgame perfect equilibrium is to play  $X_T$  as long as everyone has played on the equilibrium path. If a player deviates from playing (c, c) at some period with t rounds remaining, we play  $X_t$  as punishment. If a player deviates from playing (b, b) or (a, a), we continue on  $X_T$ .

A player that deviates with t rounds remaining gets payoff 1 in that round, and then plays the  $X_t$  subgame perfect equilibrium with payoff t+1, for a total of t+2. This will never exceed T+1, so on histories on the equilbrium path there are no profitable deviations. There are no deviations after histories when we are on  $X_t$  because they are subgame perfect equilbria, by the inductive hypothesis. Problem 2

a. This is never SPE. Player 2 has payoff 0 in equilibrium, so he can always deivate to R for payoff 1.

b. This is also never SPE, for the same reason.

c. This is always a SPE. In every period, players are playing a stage game nash equilibrium, so the strategy is subgame perfect equilibrium.

Problem 3

(a) Suppose for each cycle, (C,C) is played a times and (D,D) is played b times. Then average payoff for the cycle is  $\frac{5a+b}{a+b}$ . To make  $1.1 < \frac{5a+b}{a+b} < 1.2$ , we need 19a < b < 39a. Let's choose a = 1, b = 20. The strategy profile is for each player, play D for t = 21k + i, for  $i = 1, 2, \ldots, 20$  and play C for t = 21k if no deviation has occurred. If any deviation has occurred, play D forever.

No player has incentive to deviate when some player has deviated since (D,D) is NE of the stage game. When a player is supposed to play D at t = 21k + i, to prevent deviation we need

$$6+1\cdot\frac{\delta}{1-\delta} \leq 1\cdot\frac{1}{1-\delta} + 4\delta^{21-i}\frac{1}{1-\delta^{21}}$$

note that the right side of inequality is minimized at i = 1, so we only need to check that case. For  $\delta = 0.999$ , it holds.

When a player is supposed to play C, to prevent deviation we need

$$6 + 1 \cdot \frac{\delta}{1 - \delta} \le 1 \cdot \frac{1}{1 - \delta} + 4 \cdot \frac{1}{1 - \delta^{21}}$$

For  $\delta = 0.999$ , it holds.

(b) The strategy profile is that player 1 plays D for t = 4k + i, for i = 0, 1, 2and plays C for t = 4k + 3 and player 2 plays C for all t if no deviation has occurred. If any deviation has occured, play D forever. When (D,C) is supposed to be played, player 1 has no incentive to deviate as he gets the maximum possible payoff. For player 2, we only need to check t = 4k case (similar logic from part a) as if she were to deviate she would have maximum incentive at that case. To prevent deviation we need

$$1\cdot \frac{1}{1-\delta} \leq \delta^3 \cdot 5 \cdot \frac{1}{1-\delta^3}$$

for  $\delta = 0.999$ , it holds. When (C,C) is supposed to be played, for player 1 we need  $6 + 1 \cdot \frac{\delta}{1-\delta} \leq 6 \cdot \frac{1}{1-\delta} - \frac{1}{1-\delta^3}$  and for player 2 we need  $6 + 1 \cdot \frac{\delta}{1-\delta} \leq 5 \cdot \frac{1}{1-\delta^3}$ . Both holds for  $\delta = 0.999$ .

(c) The answer is no. To give player 1 the average payoff of more than 5.8. we have to give player to the average payoff of less than 1. Since player 2 can get at least 1 by deviation and can get at least 1 in all static NE, we cannot construct SPE where player 2 gets less than 1 on average.

No player has incentive to deviate when some player has deviated since (D,D)is NE of the stage game.

Problem 4

(a) If  $|p_1 - p_2| < c$ , we have an interior solution: there is a "mid-point"  $x^*$ such that  $0 < x^* < 1$  and kid at  $x^*$  is indifferent. In other words,

$$cx^* + p_1 = c\left(1 - x^*\right) + p_2$$

so  $x^* = \frac{1}{2} + \frac{p_2 - p_1}{2c}$ . If  $|p_1 - p_2| \ge c$ , then all kids go to one firm. To start, we find one stage (static) NE. If  $|p_1 - p_2| \ge c$ , it cannot be an equilibrium as higher price firm makes zero and has incentive to cut its price so that it can make positive profit. For  $|p_1 - p_2| < c$ , firm 1 solves

$$max_{p_1}p_1\left\{\frac{1}{2} + \frac{p_2 - p_1}{2c}\right\}$$

Taking FOC, you get  $p_1^{BR}(p_2) = \frac{c+p_2}{2}$ . Similirly,  $p_2^{BR}(p_1) = \frac{c+p_1}{2}$ . Thus,  $p_1 = p_2 = c$  as NE.

Since this NE is a unique SPE for the stage game, playing this NE for all period is a unique SPE for finite games.

(b-1) Check what would be the best response if the other firm plays  $p^*$ . If  $p^* - c \geq \frac{p^* + c}{2}$  (or  $p^* \geq 3c$ ), then BR is to charge  $p^* - c$  and capture the whole market. In this case, to make sure that there is no incentive to deviate, we need

$$\frac{p^*}{2} \left( 1 + \delta + \delta^2 + \cdots \right) \geq (p^* - c) + \frac{c}{2} \left( \delta + \delta^2 + \cdots \right)$$

$$\frac{p^*}{2} \geq (p^* - c) (1 - \delta) + \frac{c\delta}{2}$$
$$(2\delta - 1) p^* \geq (3\delta - 2) c$$

Note that  $\hat{p} = c$  here. If  $\delta > \frac{1}{2}$ , then  $p^* \ge \frac{3\delta-2}{2\delta-1}c$ . Thus, maximum  $p^*$  is  $\bar{p}$ . If  $\delta = \frac{1}{2}$ , then  $0 \ge -2c$  works for all  $p^*$ , so maximum  $p^*$  is  $\bar{p}$ .' If  $\delta < \frac{1}{2}$ , then  $p^* \le \frac{2-3\delta}{1-2\delta}c = p_{max}$ . If  $\delta \ge \frac{1}{3}$ ,  $p_{max} \ge 3c$  so maximum  $p^*$  is  $\bar{p}$ . If  $\delta < \frac{1}{3}$ ,  $p_{max} < 3c$  so the best responce is  $\frac{p^*+c}{2}$ . In this case, to make sure that there is no incentive to deviate, we need

$$\frac{p^*}{2} \quad 1 + \delta + \delta^2 + \dots \geq \frac{(p^* + c)^2}{8c} + \frac{c}{2} \quad \delta + \delta^2 + \dots$$
$$\frac{p^*}{2} \geq \frac{(p^* + c)^2}{8c} (1 - \delta) + \frac{c\delta}{2}$$

Solving, we get

$$p^* \le \frac{1+3\delta}{1-\delta}c$$

(b-2) Suppose the firm makes  $u_0$  by deviating during the war period. Note that  $u_0$  is zero if  $\hat{p} \leq 0$  and positive if  $\hat{p} > 0$ .

Let  $V_0$  as a sum of current and future profit at the war period. Then  $V_0 = \frac{\hat{p}}{2} + \frac{p^*}{2} \frac{\delta}{1-\delta}$  and to prevent deviation during the war state, we need

$$\frac{1}{1-\delta}V_0 \ge u_0 + \frac{\delta}{1-\delta}V_0$$

or  $V_0 \ge u_0$ . Since  $u_0 \ge 0$ , the best punishment is to choose  $V_0 = u_0 = 0$ . This could be done by choosing  $\hat{p} = -\frac{\delta}{1-\delta}p^*$ .

For collusion period, to prevent deviation we need

$$\frac{p^*}{2} \geq (1-\delta)(p^*-c) + \delta \cdot 0$$
$$\left(\frac{1}{2} - \delta\right)p^* \leq (1-\delta)c$$

Thus, if  $\delta \geq \frac{1}{2}$ , then  $p^* = \overline{p}$ . If  $\delta < \frac{1}{2}$ , then  $p^* = \min\left\{\overline{p}, \frac{(1-\delta)c}{\frac{1}{2}-\delta}\right\}$ .

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