6.004 Computation Structures Spring 2009

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<u>Problem 1.</u> Which of the following cannot be made to function with perfect reliability, assuming reliable components and connections. Explain your reasoning. Some of the specifications refer to "bounded time" which means there is a specified time interval, measured from the most recent input transition, after which the output is stable and valid.

- A. 🖈 A circuit that in unbounded time indicates which of two game show contestants pressed their button first.
- B. 🕇 A circuit that in bounded time indicates which of two game show contestants pressed their button first.
- C. A circuit that determines if button A was pressed before a specified deadline. Assume the circuit has an accurate internal signal that transitions from 0 to 1 when the deadline is reached. The output should be 1 if the button was pressed on or before the deadline, 0 if pressed after the deadline. The output should be valid and stable within a specified tPD of the A input transition.
- D. A circuit that in bounded time indicates which of two game show contestants pressed their button first if the presses were more than 0.1 second apart, otherwise the circuit lights up a "TIE" light.
- E. A circuit that in bounded time indicates that at least one button has been pressed by some contestant.
- F. 🕇 A circuit that in bounded time indicates that exactly one of the contestants has pressed their button. You can assume there are only two contestants.
- G. A circuit that has two parts: (a) a subcircuit that indicates which of two game show contestants pressed their button first, and (b) a subcircuit that in bounded time lights a "TIE" light if the (a) subcircuit hasn't produced an answer after 1 second. The "TIE" light should stay lit even if (a) makes a decision at some later point.
- H. A circuit that converts button presses from two contestants into the following two-bit output encoding. The circuit has two inputs, A and B, one for each contestant. A contestant's input transitions from 0 to 1 when she presses her button.

00 if neither contestant is pressing their button 01 if contestant A is pressing her button 10 if contestant B is pressing her button

11 if both contestants are pressing their buttons

The output should be valid and stable within a specified tPD of the most recent input transition.

<u>Problem 2.</u> A judge is routinely given 60 seconds to make yes-or-no decisions. Which of the following techniques allows him to do it with perfect reliability, in light of the unrealizability of a perfect bounded-time arbiter?

- A. \star He can do it reliably simply by flipping a coin.
- B. ★ He can do it reliably by asking his clerk and responding "yes" unless his clerk has answer in 55 seconds.
- C. \star He can do it realiably by always answering "no", independently of the question.

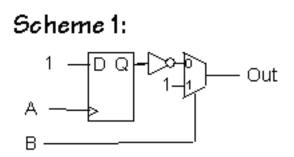
<u>Problem 3.</u> Cambridge Logic and Magic, Inc is a new startup founded by a group of ex-MIT students who dropped out of 6.004 just before Lecture 8. They are trying to develop a reliable, bounded-time arbiter that behaves as follows:

Two combinational inputs, A and B, each have a single 0-to-1 transition. The arbiter has a single output, which may be safely sampled one microsecond following the transition on the A input, at which point it is guaranteed (with 100% reliability) to have a valid logic level. If the A transition precedes the B transition by more than 100ns, the output is to be 0 when sampled; if the B transition precedes A by more than 100ns, the output is to be 1 when sampled. If the transitions are within 100ns of each other, either a 1 or a 0 may be present. In all cases, the output is stable and valid for a 10ns interval surrounding the time at which it is to be sampled.

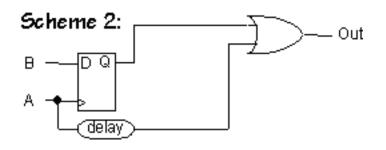
The four CLMI founders each have a proposed design for such an arbiter. Your job is to identify the problem in each design, citing a specific scenario in which it fails to meet its spec. [If you're unable to convince yourself that any scheme doesn't work, perhaps you should consider dropping out and starting a company!]

Assume in the following designs that (1) All registers start in the 0 state; (2) that transitions on each of the input signals are instantaneous and surrounded by long setup and hold times; and (3) that component timings are short compared to the one microsecond delay.

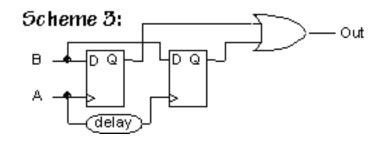
A. **★** The first scheme involves a D register and a lenient MUX. It's proponent argues that it can't go metastable, since the dynamic discipline is always obeyed on the register inputs.



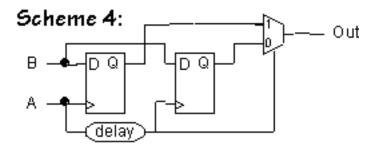
B. The second scheme involves a 50ns delay and a lenient OR gate. If the output of the register is metastable, the inventor argues, the gate's output will be forced to a valid level by the delayed copy of A. The delay is 50 ns.



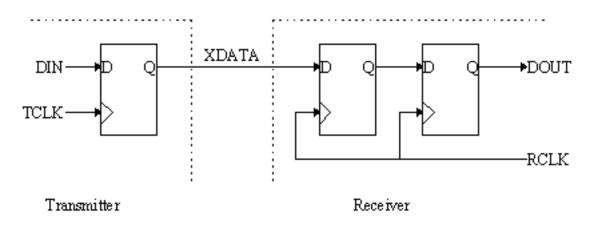
C. The third scheme involves a pair of D registers sharing the B input. The delay is 50 ns, a time chosen to ensure that if setup/hold times are violated for either of the registers the other will see valid input timing. The OR gate is lenient.



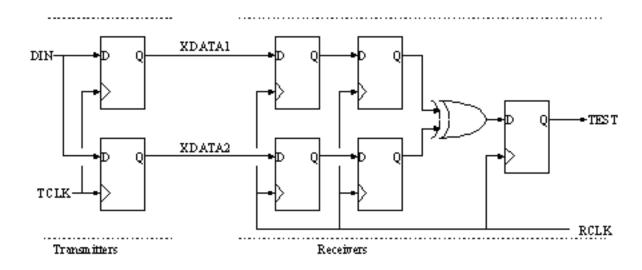
D. The fourth scheme involves a pair of D register and a delay similar to Scheme 3, but uses a MUX to select the output from a non-metastable flipflop:



<u>Problem 4.</u> Ben Bitdiddle has been hired as a consultant by Two Bit Electronics, a firm specializing in communications channels that are two bits wide. TBE has discovered that its latest product has been failing now and then at various customer sites and is worried that the design may have a problem with metastability. Ben starts by checking the schematics for one bit of the transmission channel:



- A. ★ A TBE engineer explains that although TCLK and RCLK have the same frequency, their phases may differ. Ben notes with approval the 2 D-registers in the receiver, which are configured as a **pulse synchronizer**, a circuit intended to ensure that the probability of DOUT entering the metastable state is very small. Briefly explain how the pulse synchronizer accomplishes this goal. Hint: Because the input to the receiver is asynchronous with RCLK, the left-most register in the receiver can enter a metastable state. But the probability that the input to the right-most register in the receiver does not meet that register's setup time can be made as small as desired by increasing the period of RCLK. Why?
- B. The engineer now shows Ben the schematic for a test jig that he says illustrates the problem that's been discovered. The test jig consists of two one-bit channels running in parallel with the two DIN inputs of the transmitter connected to a single data stream. The DOUT outputs of the two receivers are connected to an XOR gate that, in theory, should always produce a "0" since both receivers are getting the same data stream. In order to filter out any glitches in the XOR output introduced by small differences in tPD of the output registers of the two receivers another register has been added to the output of the XOR gate:



Every once in a while TEST is observed to be "1" for one period of RCLK. Give a brief explanation of why this happens. You can assume that DIN is meeting the setup time of the transmitter registers and that the value for DIN changes much less frequently than TCLK/RCLK. Hint: it has nothing to do with the register that was added to deglitch the TEST output. Think about what's happening in the two pulse synchronizers and why they might generate different outputs for a single clock cycle.

- C. TBE now realizes that their plan to send two bits at a time by using two independent synchronizers won't work. "Okay," they say, "we'll just transmit one bit at a time, use a single synchronizer and crank up the clock frequency." Ben observes that this circuit still won't result in completely accurate communication, i.e., the waveform produced by the receiver on DOUT still won't be the same as the waveform produced by the transmitter on XDATA. Give a brief explanation of why Ben is right.
- D. Ben explains that communications between asynchronous systems is a tricky business! If the input stream isn't too fast compared to the receive clock, a pulse synchronizer will reliably reproduce any transitions in the input waveform, although the exact timing between the transitions in the output waveform will not the same as in the incoming waveform.

[Hard] Can you think of a way to encode a bit stream using transitions rather than levels to convey the information? Describe your encoding scheme and the signals that would be output by your receiver.

Problem 5. From the weekly Campus Police report: Ben Bitdiddle left the Bit Bucket Cafe and

approached the nearby fork in the road. He hits the barrier in the middle of fork, later explaining "I can't be expected to decide which side of the fork to take in bounded time!."

A. Is the accident Ben's fault?

Problem 6. Here's an everyday situation: you have a traffic intersection controlled by a traffic light with

the normal green, yellow, red sequence. There's a 55 MPH speed limit and the yellow and green periods of the traffic light are sufficiently long. I.e., the yellow is long enough to ensure that a car traveling at or below the speed limit can come to a halt before the stop line, and the green is long enough to ensure that a car stopped at the stop line can accelerate and cross the intersection before the green period is over.

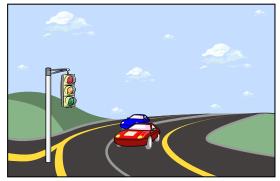


Image by MIT OpenCourseWare.

The car is supplied with digital GREEN, YELLOW and RED inputs from the traffic light and also has an analog POSITION input giving the distance from the stop line. The car's controller produces a digital GO output that controls the motion of the car. The car is is expected to behave reasonably -- e.g. to make normal progress in the absence of a red light.

- A. Is it possible to build a controller that reliably obeys Law #1: "don't cross the stop line when the light is red"?
- B. Is it possible to build a controller that reliably obeys Law #2: "don't be in the intersection when the light is red"?