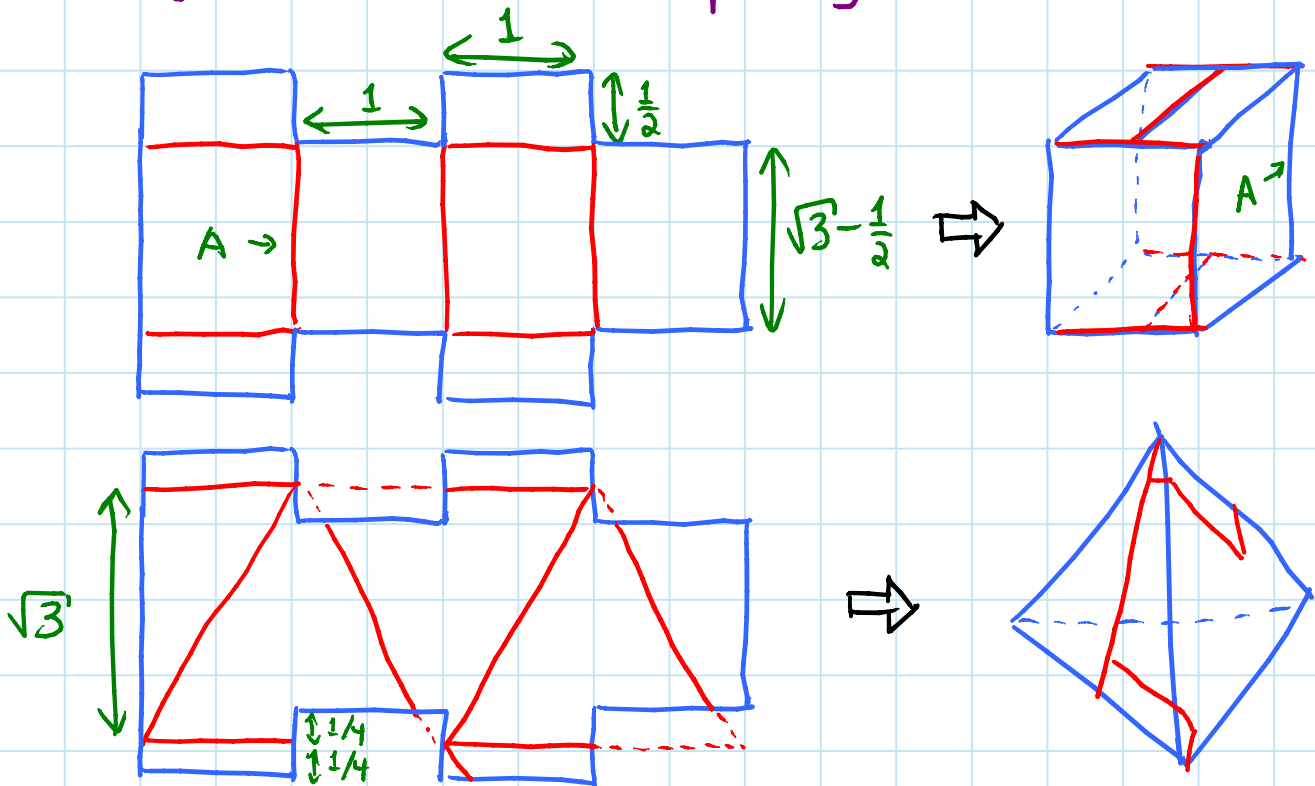


## Polyhedron unfold/refold "dissection": (common unfoldings)

**OPEN**: which polyhedra  $P, Q$  are connected by  $P \rightarrow$  (general) unfolding  $\rightarrow$  (general) gluing  $\rightarrow Q$ ?  
 or a sequence of unfolding, gluing, ...?  
 - true for any two Platonics? [M. Demaine, 1998]  
 e.g. regular tetrahedron & cube

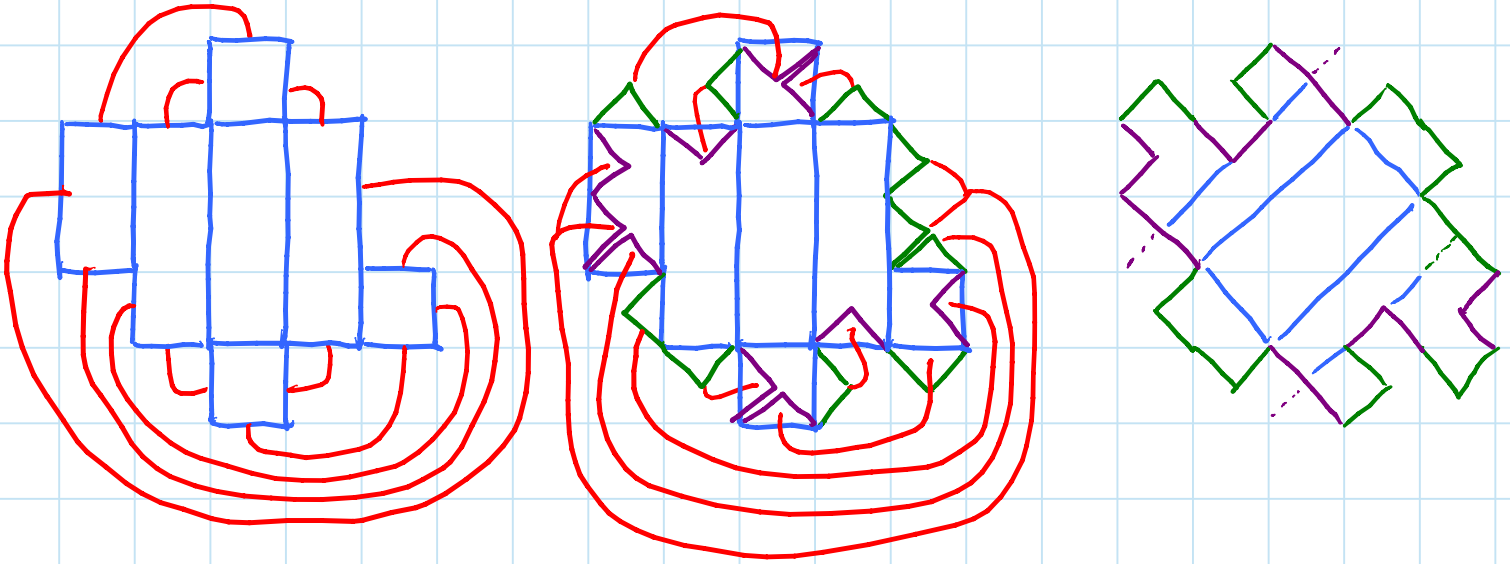
### Interesting examples:

- regular tetrahedron  $\rightarrow 1 \times 1 \times \underbrace{1.232}_{\sqrt{3} - \frac{1}{2}}$  box possible  
 [Hirata 2000; GFALOP p.425]



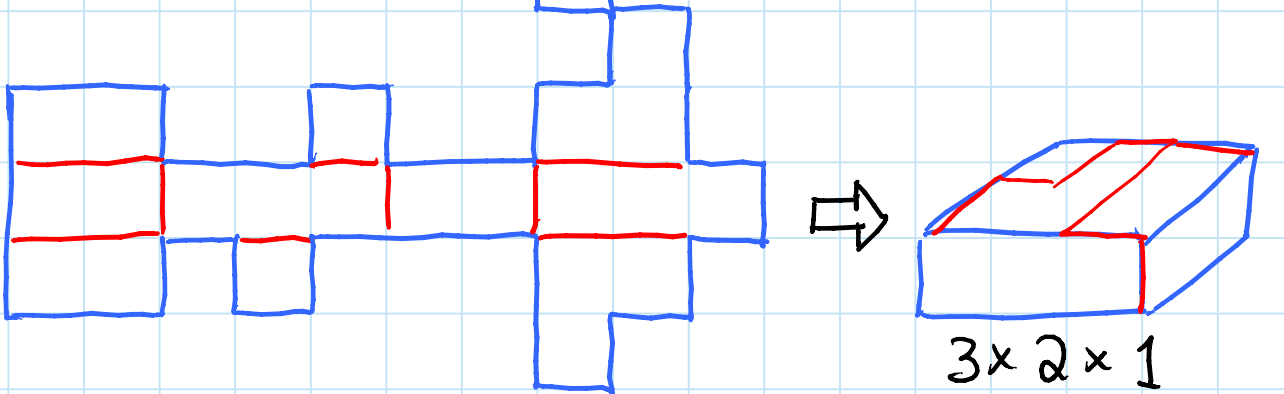
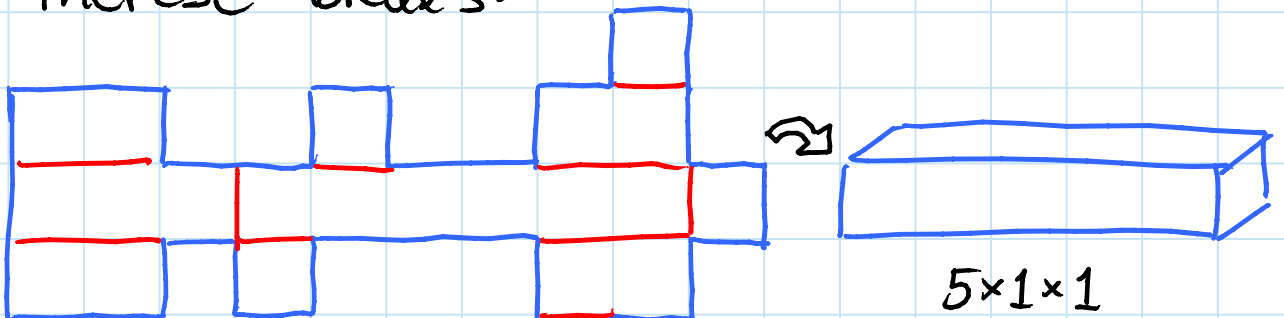
- regular (icosahedron) & tetrahedron  
 (octahedron) with equal faces [Horiyama & Uehara 2010]  
 (cube)
- conjectured fractal common unfolding of  
 regular tetrahedron & cube [Shirakawa 2010]

Box  $\rightarrow$  box: [Biedl, Chan, Demaine, Demaine, Lubiw, Munro, Shallit 1999]  
 - Timothy Chan's solution: [GFALOP p. 425]



$\sqrt{2} \times \sqrt{2} \times 3\sqrt{2} \rightarrow$  add tabs  $\rightarrow$   $1 \times 2 \times 4$  box

- Therese Biedl's:



- also  $8 \times 1 \times 1 \rightarrow 5 \times 2 \times 1$

- OPEN: when is  $a \times b \times c \rightarrow d \times e \times f$  possible?

## More boxes: [Uehara - CCCG 2008]


- random generation algorithm for orthogonal grid unfoldings (like Biedl's above)
- found  $> 25,000$  so far!
- $1 \times 1 \times (6k+2) \leftrightarrow 1 \times 5 \times 2k$  &  
 $1 \times 1 \times (8k+11) \leftrightarrow 1 \times 3 \times (4k+5)$  always possible  
 $\Rightarrow$  infinitely many
- tiling example

OPEN: common unfolding of 3 boxes?

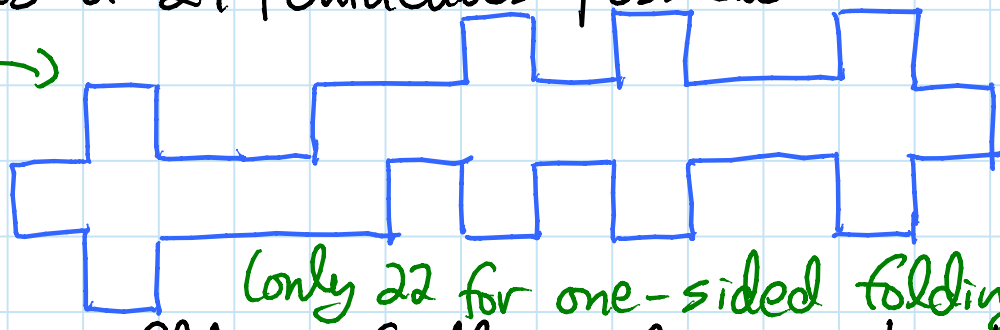
Cubigami: [Knuth & Miller 2005; Recent Toys]  
common unfolding of all tree tetracubes

common  
surface  
→ area

Generalization: [Aloupis, Benbernou, Bose, Collette, Demaine,  
Demaine, Duić, Dujmović, Iacono, Langerman, Morin 2010]

- no common unfolding of all tree pentacubes (trying all 1,099,511,627,776 unfoldings of )
- but 23 of 24 pentacubes possible

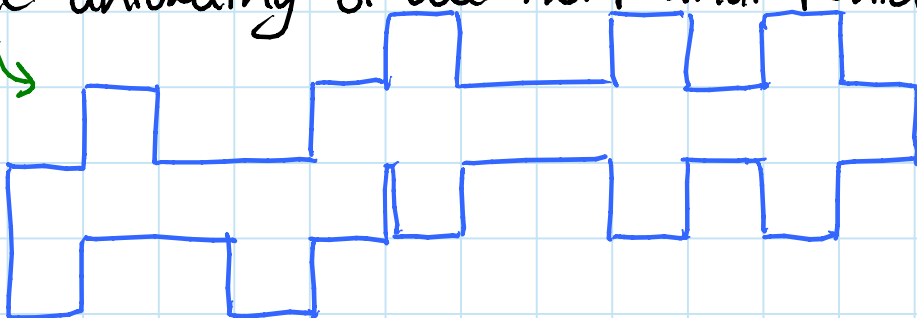
e.g. →



only 22 for one-sided foldings)

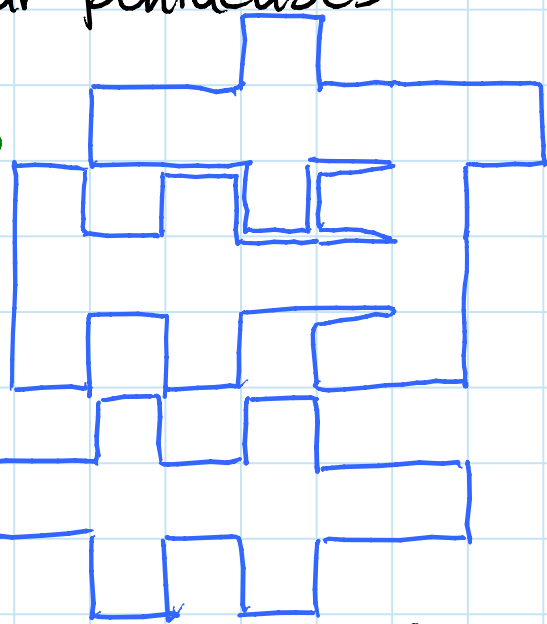
- unique unfolding of all nonplanar pentacubes (& 22 total)

→



- 492 unfoldings of all planar pentacubes (none fold to any nonplanar pentacube) →

- no common unfolding of all path pentacubes
- planar hexacubes! →



**OPEN**: 7?

- **OPEN**: any two polycubes without common unf.?


# Orthogonal polygons vs. polyhedra:

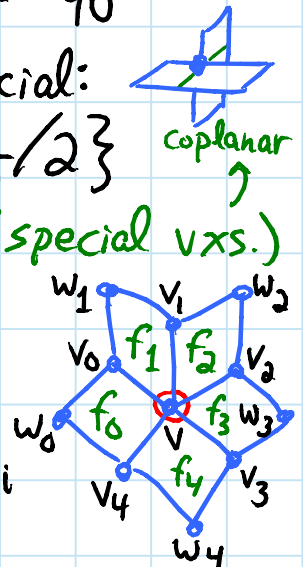
can an orthogonal polygon with orthogonal creases fold into nonorthogonal polyhedra?

- answer depends on allowed genus

<u>Genus</u>	<u>NonOrthog.?</u>	<u>Reference</u>
$\emptyset$	NO	[Donoso & O'Rourke 2002]
1 & 2	NO	[Biedl et al. 2002]
3 - 5	<b>OPEN</b>	
6	YES	
$\geq 7$	YES	

## Proof sketch that genus $\leq 2 \Rightarrow$ impossible:


- edge green if dihedral multiple of  $90^\circ$ ; red else
- focus on red subgraph; coalesce  $\xrightarrow{\text{deg. 2}} \text{---}$
- faces not planar, but face angles still mult. of  $90^\circ$  & dihedral angles still not mult. of  $90^\circ$
- vertices have degree  $\geq 4$ ; deg. 4 special: 
- claim:  $V \leq 8(g-1) - \max\{D_{\neq 5}, F_{\geq 5}/2\}$  coplanar  $\uparrow$   
(proof uses Euler's Theorem + analysis of special vxs.)
- degree  $\geq 4$  &  $V \geq 1 \Rightarrow V \geq 5$
- claim  $\Rightarrow D_{\neq 5} \leq 3 \Rightarrow D_5 \geq 2 \rightarrow V$
- $V \geq 6 \Rightarrow F_{\geq 5} \leq 4 \Rightarrow$  some  $f_i$  degree 4  $\rightarrow w_i$
- $V \geq 7 \Rightarrow F_{\geq 5} \leq 2 \Rightarrow \dots \Rightarrow$  all  $f_i$  degree 4
- $w_i \neq w_{i+1} \Rightarrow \geq 3$  distinct  $\Rightarrow V \geq 9$  but  $V \leq 8 \quad \square$



## Smooth Alexandrov: [Pogorelov 1973]

- every convex metric, topologically a sphere, is realized by a unique convex surface, possibly degenerating to flat doubly covered convex shape
- proof idea: take limits of polyhedral approximation

## D-forms: [Tony Wills; Potmann & Wallner 2001]

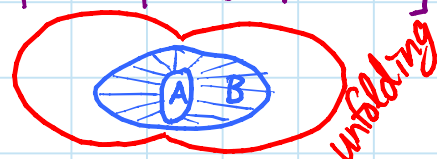
- take two convex smooth shapes of the same length 
- identify two boundary points of two shapes
- glue around from there
- smooth Alexandrov  $\Rightarrow$  get convex surface

## Results: [Demaine & Price 2009]

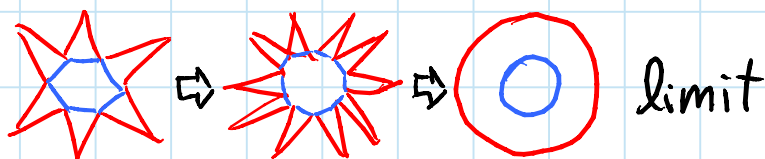
- no creases except along seam
- form = convex hull of its seam

# Unfolding smooth prisms: [Benbernou, Cahn, O'Rourke 2004]

= convex hull of two parallel smooth convex shapes



- keep bottom B, unroll side ribs, place top A  
*volcano unfolding*      *nontrivial*
- simple example: pyramid (A = single point)



- surface area not preserved (unlike regular unfolding)
- like source unfolding of sphere (*below*)

- always works, yet:

**OPEN**: does every (discrete) prismatoid have an edge unfolding?

↑  
from Lecture 13

## Wrapping smooth surfaces with flat paper:

[Demaine, Demaine, Iacono, Langerman 2007]

- impossible with finite number of creases to make all points have nonzero curvature
- idea: allow folding to shrink some intrinsic distances on paper (contractive mapping)
  - simulate by crinkling paper:

## Burago & Zalgaller Theorem: [1996]

any contractive  $C^2$ -immersion of a polygon (or more generally, a polyhedral metric) admits a  $C^0$ -approximation (within  $\varepsilon \forall \varepsilon > 0$ ) by isometric piecewise-linear  $C^0$ -immersions

- also, noncrossingness is preserved

Stretched path = isometrically folded (unshrunk) path

- optimal wrapping should have one; else scale

Stretched wrapping = stretched path between every two points

## Source wrapping of convex surfaces:



- stretched paths along all shortest paths from  $x$
- cut at ridge tree / cut locus
- e.g. unit sphere  $\Rightarrow$  disk of radius  $\pi$ , area  $\pi^2$ , perimeter  $2\pi$

## Strip wrapping:

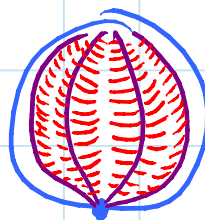
- e.g. unit sphere  $\Rightarrow$  area  $\rightarrow 4\pi$ , perimeter  $\rightarrow \infty$



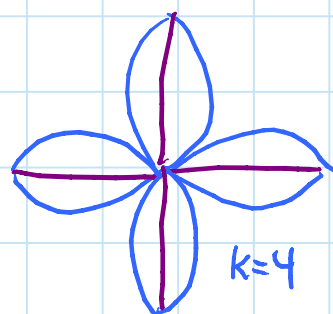
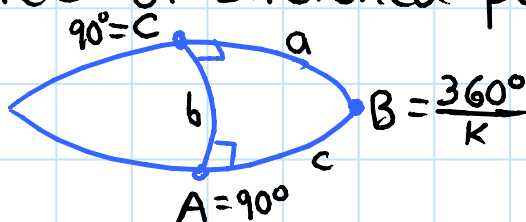
# Wrapping smooth surfaces with flat paper: (cont'd)

## Petal wrapping of sphere:

- $k$  stretched paths from south to north pole
- perpendicular stretched paths till meet another  $\Rightarrow 360^\circ/k$  "orange peel" for each primary path
- depth-2 tree of stretched paths



Voronoi diagram



- Spherical Law of Cosines  $\Rightarrow$

$$\cos C = \underbrace{-\cos A}_{\emptyset} \cos B + \underbrace{\sin A}_{1} \sin B \cos c$$

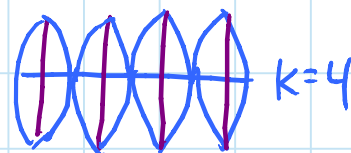
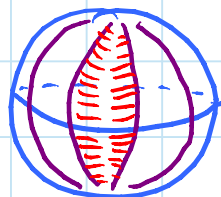
$$\Rightarrow \sin C = \sqrt{1 - \sin^2 B \cos^2 c}$$

- Spherical Law of Sines  $\Rightarrow \frac{\sin B}{\sin b} = \frac{\sin C}{\sin c}$

$$\Rightarrow b = b(c) = \arcsin \frac{\sin c}{\sqrt{\frac{1}{\sin^2 B} - \cos^2 c}}$$

## Comb wrapping: depth-3 tree

- stretched path around equator
- $k$  paths from equator to north pole, ditto for south pole
- perpendicular stretched paths
- same petals, different gluing



## Wrapping smooth surfaces with flat paper: (cont'd)

Both petal & comb wrappings: area  $\rightarrow 4\pi$   
perimeter  $\rightarrow \infty$

Real Mozartkugel wrappings:

- square containing 4-petal [Fürst]
  - $\sqrt{2}\pi \times \sqrt{2}\pi$
  - $\Rightarrow$  area  $2\pi^2$ , perimeter  $8\pi/\sqrt{2} \approx 5.7\pi$
- rectangle containing 4-comb [Mirabell]  $\nearrow$ 
  - $\pi \times 2\pi$
  - $\Rightarrow$  area  $2\pi^2$  (!), perimeter  $6\pi >$

Better Mozartkugel wrappings:

- equilateral triangle containing 3-petal has area  $1.9983\pi^2$  ( $\approx 0.1\%$  improvement)
- packing 3-petals  $\Rightarrow 1.6033\pi^2$  area each
- packing k-combs  $\Rightarrow 1.3333\pi^2$  area each
- vs. optimal:  $4\pi = 1.2732\pi^2$

**OPEN**: what is the best area for given perimeter?  
(Pareto curve)

**OPEN**: what is the minimum possible perimeter?

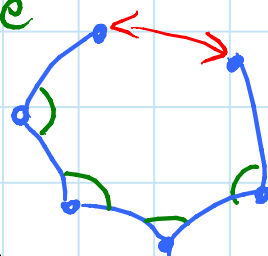
**OPEN**: what about smooth surfaces other than the sphere?

# How to prove contractiveness of these wrappings?

## Cauchy's Arm Lemma on a growing sphere:

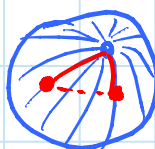
[Abel, Charlton, Collette, Demaine, Demaine, Langerman, O'Rourke, Pinciu, Toussaint 2008] e.g. plane

open convex chain on a sphere, redrawing on a larger-radius sphere with matching edge lengths & angles, increases the length of the closing edge



## Contractive corollaries:

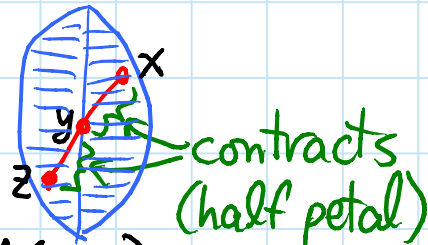
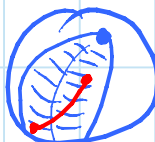
- source wrapping



- half petal



- petal

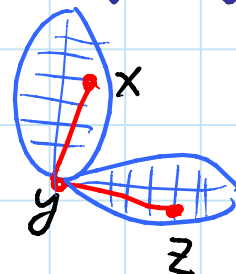


- plane:  $d(x, z) = d(x, y) + d(y, z)$

- sphere:  $\geq d'(x, y) + d'(y, z)$

$\geq d'(x, z)$  ( $\Delta$  inequality)

- petal unfolding: same argument (shortest path between x & z)



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