Transportation Management Vehicle Routing

Chris Caplice ESD.260/15.770/1.260 Logistics Systems Nov 2006

Local Routing

Large Number of Network Problems – we will look at four

Shortest Path Problem

- Given: One origin, one destination
- Find: Shortest path from single origin to single destination
- Transportation Problem
 - Given: Many origins, many destinations, constrained supply
 - Find: Flow from origins to destinations
- Traveling Salesman Problem
 - Given: One origin, many destinations, sequential stops, one vehicle
 - Find: Shortest path connecting each stop once and only once
- Vehicle routing Problem
 - Given: One origin, many destinations, many capacitated vehicles
 - Find: Lowest cost tours of vehicles to destinations

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Shortest Path Problem

- Find the shortest path in a network between two nodes – or from one node to all others
- Result is used as base for other analysis
- Connects physical to operational network
- Issues
 - What route in practice is used? Shortest? Fastest? Unrestricted?
 - Frequency of updating the network
 - Using time versus distance (triangle inequality)
 - Impact of real-time changes in congestion
 - Speed of calculating versus look-up

Shortest Path



Shortest Path Matrix

i\j	1	2	3	4	• • •	n
1		d ₁₂	d ₁₃	d ₁₄		d _{1n}
2			d ₂₃	d ₂₄		d _{2n}
3				d ₃₄		d _{3n}
4						d _{4n}
• • •						
n			*****			

Network

- Arc/Link & Nodes
- Cost is on nodes, c_{ij}
- Think of a string model
- Basic SP Algorithm (s to t)
 - 1. Start at origin node, s=i
 - 2. Label each adjacent nodes, j, $L'_j = L_i + C_{ij}$ iff $L'_j < L_j$
 - 3. Pick node with lowest label, set it to i, go to step 2
 - 4. Stop when you hit node t
- Building Shortest Path Tree
- Many, many variations on this algorithm,
 - Label Setting
 - Label Correcting

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Transportation Problem



Traveling Salesman Problem

- Starting from an origin, find the minimum distance required to visit each destination once and only once and return to origin.
 - *m*-TSP: best tour for *m* salesmen
- Very old problem ~1832
 - For history, see: <u>http://www.tsp.gatech.edu/index.html</u>



TSP Solution Approaches

- Heuristics
 - Construction
 - Nearest neighbor
 - Greedy (complete graph, pick shortest edge until Hamiltonian path)
 - Sweep (example of Cluster-First, Route-Second)
 - Space filling curve (example of Route-First, Cluster-Second)
 - Insertion (nearest, cheapest)
 - Savings (Clarke-Wright)
 - Local Improvement
 - 2-opt
 - 3-opt
 - Meta-heuristics
 - Tabu Search
 - Ant System
 - Simulated Annealing
 - Genetic Algorithms
 - Constraint Programming

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Traveling Salesman Problem

Nearest Neighbor Heuristic

- Start at any node and connect tour to closest adjacent node
- In practice 20% above optimal
- Insertion Heuristic
 - Form some sub tour (convex hull) and add in the nearest/furthest/cheapest/random node one at a time
 - In practice 19% / 9% / 16% / 11% above optimal
- 2-Opt Heuristic
 - Method of improving a solution
 - Select two edges (a,b) and (c,d) where total tour distance decreases the most if reformed as (a,c) and (b,d).

Vehicle Routing Problem

- Find minimum cost tours from single origin to multiple destinations using multiple vehicles
- Who needs to solve the problem?
 - Shippers retailers, distributors, manufacturers
 - Carriers LTL, package
 - Service companies repair, waste, utility, postal, snow removal
- Types of problems
 - Commercial delivery (retailers, distributors, manufacturers)
 - Commercial pickup (retailers, distributors, manufacturers)
 - Mixed pickup & delivery (LTL and package carriers)
 - Residential appointment (online grocery, medical gases, repair)
 - Residential sweep (postal, waste, utility, snow removal)

Initial Routes



Optimized Routes



VRP is NP-Hard

Difficult to evaluate quality by enumeration

Combinatorial Growth

		Cı	ustomers			
		total	on the route	Ways to select customers for the route	Ways to select and sequence the route	Hours of work to evaluate one per second
		10	3	120	720	0.20
		20	3	1,140	6,840	1.9
- V	3 Stops	30	3	4,060	24,360	6.8
		40	3	9,880	59,280	16
		50	3	19,600	117,600	33
		60	3	34,220	205,320	57
		70	3	54,740	328,440	91
		80	3	82,160	492,960	137
		90	3	117,480	704,880	196
		100	3	161,700	970,200	270
						Dave of work to avaluate and per second
		10	5	252	30.240	Days of vork to evaluate one per second
		20	5	15 504	1 860 480	22
		30	5	142 506	17 100 720	198
		40	5	658,008	78 960 960	914
_		50	5	2 118 760	254 251 200	2 943
	5 stops	60	5	5 461 512	655,381,440	7.585
	0 00000	70	5	12 103 014	1 452 361 680	16.810
		80	5	24.040.016	2.884.801.920	33,389
		90	5	43,949,268	5,273,912,160	61.041
		100	5	75,287,520	9,034,502,400	104,566
		total	on the route			Years of work to evaluate one per second
		10	10	1	3,628,800	0.12
		20	10	184,756	670,442,572,800	21,260
	10 ctopc	30	10	30,045,015	109,027,350,432,000	3,457,235
\sim	10 51005	40	10	847,660,528	3,075,990,524,006,400	97,539,020
		50	10	10,272,278,170	37,276,043,023,296,000	1,182,015,570
		60	10	75,394,027,566	273,589,847,231,501,000	8,675,477,145
		70	10	396,704,524,216	1,439,561,377,475,020,000	45,648,191,828
		80	10	1,646,492,110,120	5,974,790,569,203,460,000	189,459,366,096
		90	10	5,720,645,481,903	20,759,078,324,729,600,000	658,266,055,452
		100	10	17,310,309,456,440	62,815,650,955,529,500,000	1,991,871,225,125
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Vehicle Routing Problems

General Approaches

- Heuristics
 - Route first Cluster second
 - Space filling curve
 - Any earlier heuristic can be used
 - Cluster first Route second
 - Sweep Algorithm
 - Savings (Clarke-Wright)
- Optimal
 - MILP Column Generation



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Savings Algorithm

Clark-Wright Algorithm

- Serve each node directly
- Identify savings for combining two nodes on same tour
- Add nodes together if savings >0

•
$$2C_{0i} + 2C_{0j} > C_{0i} + C_{ij} + C_{j0}$$

• Savings =
$$c_{0i} + c_{j0} - c_{ij}$$

Shortest Path Matrix

i\j	0	1	2	3	4
0		10	15	19	22
1			8	23	35
2				12	21
3					5



Suppose Max Capacity = 3 Suppose Max Capacity = 3 Savings = $c_{0i} + c_{j0} - c_{ij}$ S(1,2) = 10 + 15 - 8 = 17 S(1,3) = 10 + 19 - 23 = 6S(1,4) = 10 + 22 - 35 = -3

•
$$S(2,3) = 15 + 19 - 12 = 22$$

•
$$S(2,4) = 15 + 22 - 21 = 16$$

•
$$S(3,4) = 19 + 22 - 5 = 36$$

Shortest Path Matrix

i\j	0	1	2	3	4
0		10	15	19	22
1			8	23	35
2				12	21
3					5



Benefits of this approach?

Optimal Approach – MILP w/CG

	Route 1	Route 2	Route 3			Route M	
	C1	C2	C3			Cm	
Stop A	1	0	0	1	0	1	1
Stop B	1	1	0	0	1	0	1
Stop C	1	1	1	1	0	0	1
Stop D	0	1	1	0	1	1	1
Stop E	0	0	1	1	0	0	1
Stop F	0	0	0	0	1	0	1
Stop G	0	0	0	0	0	1	1
	0	0	0	0	0	0	1
	0	0	0	0	0	0	1
Stop N	0	0	0	0	0	0	1

Each Row is a stop Each Column is a generated vehicle route and its cost Each matrix coefficient, a_{ij} , is {0,1}, identifying the stops on the j'th route Define Y_j, {0,1}, "1" if the route is used ; else "0"

Subject to:

Minimize: $\sum_{i} C_{j} Y_{j}$

 $\sum_{j=1}^{j} a_{ij} Y_j \ge D_i \text{ ; for all I}$

 $Y_i = \{0,1\}$, for all J

Same Example

Shortest Path Matrix

												i\j	j 0		1	2	3	4
Each tour is a column											0			10	15	19	22	
 How are tours generated? Could each column be a solution? How could this be enhanced? 												1				8	23	35
												2					12	21
– 110 W	cou	ind t	1113 1				, u					3						5
Total Dist																		
132																		
															-			
Dec Var	1	1	1	1	0	0	0	0	0	0	0	0	0	0				
	٢	2	3	4	5	6	7	8	9	10	11	12	13	14				
	ute	ute	ute	ute	ute	ute	ute	ute	ute	ute	ute	ute	ute	ute				
	Ro	Ro	Ro	Ro	Ro	Ro	Ro	Ro	Ro	Ro	Ro	Ro	Ro	Ro	Su	m	RH	S
Stop 1	1	0	0	0	1	1	1	0	0	0	1	1	1	0	1	>=	1	
Stop 2	0	1	0	0	1	0	0	1	1	0	1	1	0	1	1	>=	1	
Stop 3	0	0	1	0	0	1	0	1	0	1	1	0	1	1	1	>=	-1	
Stop 4	0	0	0	1	0	0	1	0	1	1	0	1	1	1	1	>=	1	
Capacity	1	1	1	1	2	2	2	2	2	2	3	3	3	3	_			
Distance	20	30	38	44	33	52	67	46	58	46	59	61	60	54				

Regardless of Approach

Rules of Thumb

- Good routes are "rounded", not "star shaped"
- Good routes don't cross themselves or others
- Good sectors are "pie shaped", not "checker board"
- Good solutions "look like a daisy"
- Good Practice Tips
 - Always use a Preview-Analyze-Review methodology
 - Periodically visit the internal logic within the TMS
 - Never discount the salty expert who has been doing this longer than you've been alive
 - Identify all special conditions (customer A must be delivered to first) and then validate or reject them

Other Extensions to VRP

More dimensions/elements

- Sourcing
 - Multiple depot
 - Dynamic sourcing (depot varies)
 - Order
 - Multiple dimensions (e.g. cube, weight)
 - Mixed pickup and delivery
 - Time window
 - "Vendor Managed Inventory"
- Plan
 - Fixed / Static / Master
 - Variable / Dynamic / Daily
 - Zone
 - Real-time dispatch
 - Resource
 - Backhaul
 - Continuous moves

Academic problems

- Multiple Depot VRP (MDVRP)
- Multi-commodity VRP
- Vehicle Routing Problem with Pick-up and Delivering (VRPPD)
- VRP with time windows (VRPTW)
- Inventory Routing Problem (IRP)
- Stochastic VRP (SVRP) minimize expected costs for satisfying realized demand/customers
- Dynamic VRP redirect trucks during the execution of their route to accommodate new orders
- Vehicle Routing Problem with Backhauls (VRPB)

Fixed vs. Dynamic Route Plans

- Fixed/static routes
 - Routes repeat on a cycle
 - Daily, weekly, whenever there is sufficient demand
 - Routes are changed when customer base changes
 - Quarterly, annually
 - Routes are based on "forecast" demand
 - Routes are designed for "heavy days" related to truck capacity and driver hours
 - Primary advantages
 - Driver familiarity
 - Ease of execution
 - Primary disadvantages
 - Inefficiency caused by variability
 - Difficulty of efficient customer day assignment

- Variable/dynamic routes
 - Routes change continually
 - Typically every day
 - Routes based on "actual" shipment requirements
 - Routes are designed for vehicle and driver constraints
 - Primary advantages
 - Utilization of trucks and drivers
 - Flexibility in customer ordering
 - Primary disadvantages
 - Difficulty of determining optimum routes
 - Difficulty of maintaining route planning process
 - Execution may not match plan

Real-World Issues

- The real world does not behave according to uniform assumptions
 - Dock configuration
 - Dock hours
 - Trailer types
 - Moveable bulkheads (bulk liquids, grocery reefers)
 - Truck types
 - Truck-trailer combos: doubles & triples (pups)
 - Compatibility: order-vehicle, order-order, vehicle-site
 - Preferred customers (big box)
 - Driver preferences (seniority, local knowledge)
 - Driver skills (service technician)
 - Rush hour traffic
 - Real-time dispatching (deployed vehicles)
 - Refueling
 - Maintenance

Element Interactions

Truck & Trailer

Trailers the tractor can handle – length, pups, specialized (e.g. car hauler)

Vehicle & Customer

Must be able to visit the customer (loading dock, cornering, parking)

Vehicle & Order

 Products may not be deliverable on certain resources -- HazMat, loading/handling equipment (tanks, racks), capabilities (refrigeration), physical dimensions, etc.

Vehicle & Driver

Not licensed for the truck, not able to load/unload trailer

Order & Order

Products may not mix (lumber & light bulbs, bottled water & dehydrated food, etc.)

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Manual Planning

- Plan using paper, pencil, and experience
 Advantages
 - Cheap and easy
- Challenges
 - Cannot generate multiple solutions
 - Difficult to evaluate result
 - Decentralized

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Interactive GIS

Plan using human intuition, guided by simple heuristics

Advantages

Evaluation is easier (distance, time, cost calculations, and visual)

Challenges

- Time consuming (and typically there is limited time for planning)
- Requires "super-users"
 - Need technical aptitude
 - Requires regular training
- Typically decentralized

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Automated Heuristics

- Plan using construction, local improvement, & other heuristics
- Advantages
 - Provides solutions relatively quickly
- Challenges
 - Solution quality hard to predict
 - Heuristics that work well for one problem may work poorly for another
 - Solution quality from heuristics can change drastically when the data changes
 - Hard to know when to settle on a solution
 - Complexity
 - Not as good if there are complex constraints or shipments vary in size
 - Need sophisticated expert to improve or tune
 - Typically users stick with the same approach and manually edit plans

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Optimization

Column generation and set covering IP

Advantages

 Determines best solution among the options considered

Challenges

- Quality depends on quality of options created (column generation)
- Requires significant computing power (parallel computing is advantageous)
- Requires regular maintenance by domain and technology experts

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Questions?