## Complexity of Games \& Puzzles [De maine, Hearn \&̛ many others]



## Constraint Logic [Hearn \& Demaine 2009]



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## Constraint Graphs

$$
\begin{aligned}
& \text { Machine = graph, } \\
& \text { red of } 6 \text { lue edges }
\end{aligned}
$$



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## Constraint Graphs

## Machine state <br> =orientation



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## Constraint Logic



> Rule: at least 2 units
> incoming at a vertex

Move: reverse an edge, preserving Rule
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## AND vertex



Rule: at least 2 units
incoming at a vertex

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## SPLIT vertex



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## OR vertex



Rule: at least 2 units
incoming at a vertex

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## Decision Problem



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## Decision Problem



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## Sliding -Block Puzzles

# DADSPUZZLER Few solve it 

 IT CAN BE DONE
"TAKE ONE HOME"

Courtesy of Dr. Jim Starer. Used with permission.

## Sliding-Block Puzzles [Hearn \& Demaine 2002]


(a) AND


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## Sliding-Block Puzzles [Hearn \& Demaine 2002]



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## Wiring Vertices Together



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## Red-Blue Conversion


assume an even number of conversions
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## Red-Blue Conversion


assume an even number of conversions
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Boole an Formulas


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## Quantified Boole an Formulas ( $Q \mathcal{B F}$ )

$$
\forall x \exists y \forall w \cdots \exists z[(x \vee y) \wedge \cdots \wedge(\bar{z} \vee x \vee \bar{w})]
$$



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## Existential Quantifier



## satisfied out

## Universal Quantifier



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## Latch



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## Universal Quantifier



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## Crossover Gadget



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## OR from <br> Protector OR


(b) Protected $O R$.

H


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## Rush Hour [Hearn \& Demaine 2002]


(b) AND

(c) Protected OR

PS PACE-comple teness Known[Flake \& Baum 2002]
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## Triangular Rush Hour

[Hearn \& Demaine 2009]


(a) AND vertex

(b) Connector

(c) OR vertex

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## Open: $1 \times 1$ Rush Hour [Tromp ơ Cilibrasi 2008]

- P or $\operatorname{PS} P \mathcal{A C E}$-complete or ..?


Image courtesy of John Tromp. Used with permission.

## Plank Puzzles

## [Hearn 2004]

output

(a) AND

(b) OR

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## Sokoban [Hearn $\Leftarrow$ Demaine 2002]



PS PACE-comple teness Known [Culberson 1998]
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## Constraint Logic [Yearn \& De maine 2009]

 EComputation

| Robert A. M earn |
| :--- |
| trike | Erik D. Devnaine

Games, Puzzles,

2 players (game)
team, imperfect info Copyright (2009) From Games, Puzzles, and Computation by Robert A. Hearn and Erik D. Demaine. Reproduced by permission of Taylor and Francis Group, LLC, a division of Inform plc.

## Constraint Logic [Hearn \& Demaine 2009]



0 players
(simulation)


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## Amazons

[Hearn 2005]


(a) CHOICE

(b) AND


(c) FANOUT

(d) OR

(e) VARIABLE

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## ES. 268 The Mathematics in Toys and Games

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